Embedding trees in dense graphs

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joint work with T. Klimošová and D. Piguet

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Definition (Extremal graph theory, Bollobás 1976)

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Alternative definition: substructures in graphs

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Graph G has n vertices. If G has more than $n^2/4$ edges then it contains a triangle.

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Density of edges vs. density of triangles (Razborov 2008)



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Generalisations? Other cliques (Turán 1941), asymptotically for all non-bipartite graphs (Erdős-Stone 1946) The answer for C_4 is of order $n^{3/2}$, lower bound via finite projective planes. What is the answer for trees? Fix any tree T on k vertices. There are graphs with average degree k - 2 that do not contain T.

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Conjecture (Erdős-Sós)

Any graph with average degree greater than k - 2 contains any tree on k vertices as a subgraph.

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Partial results:

- special trees (paths Erdős, Gallai 1959)
- special graphs (without C₄ Saclé, Wozniak 1997)
- n and k differ by constant (Görlich, Żak 2016)

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One can get reasonably close if the size of the tree is comparable with the size of the graph. Below Δ is maximum degree and $\overline{\deg}$ average degree.

Theorem (R. 2019), also (Besomi, Pavez-Signé, Stein 2019+)

Let \mathcal{T} be a class of trees such that $\forall T \in \mathcal{T} : \Delta(T) \in o(|T|)$. Then any graph G with $d\overline{eg}(G) = |T| + o(|G|)$ contains any $T \in \mathcal{T}$.

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Conjecture (Simonovits)

If at least rn vertices of G have degree at least k, then G contains any tree with k + 1 vertices and at most r(k + 1) vertices in one colour class as a subgraph.

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If at least *rn* vertices of *G* have degree at least *k*, then *G* contains any tree with k + 1 vertices and at most r(k + 1) vertices in one colour class as a subgraph.

Theorem (Klimošová, Piguet, R. 2019+)

If at least *rn* vertices of *G* have degree at least k+o(n), then *G* contains any tree with k + 1 vertices and at most r(k + 1) vertices in one colour class as a subgraph.

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Turns out that one can get this 'r' trade-off also in the proof of previous Erdős-Sós result.

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- Aim is to find suitable decomposition of the tree such that small subtrees are embedded by Szemerédi and the macro structure by us.
- We reduced the problem to a certain fractional variant of itself. But now we have much simpler tree structure to work with.



Proof.

Condition on the maximum degree actually gives even simpler decomposition. After decomposition of G and T look at a high degree cluster of G and a maximal matching in its neighbourhood. Provide (almost) greedy algorithm for embedding.



Proof of the Loebl-Komlós-Sós result

Proof.

After decomposition of G and T 'discharge' into several configurations, embedding for each one being straightforward.



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